

# Study on the digitization dual combinatorics and convex case

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**Abstract.** The action of a translation on a continuous object before its digitization generates several digitizations. The *dual*, introduced by the authors in a previous paper, stands for these digitizations in function of the translation parameters. This paper focuses on the combinatorics of the dual by making a link between the digitization number and the boundary curve, especially through its dual representation. The convex case is then studied and a few significant examples are exhibited.

## 1 Introduction

For a given grid step and a given digitization method, a planar object produces several digitizations in function of its position on the grid. The object digital properties and digitally estimated characteristics depend on the obtained digitization. Thereby, this study of the digitization variability is an important issue in image analysis.

This field has been explored for some geometrical primitives. For instance, the set of straight segment digitizations in function of the segment slope and offset is known as the segment *preimage* and is used for digital straight segment recognition [2]. Several papers are also dedicated to the study of the generation and combinatorics of the disc digitization set in function of its radius and center position [3, 9, 13, 5, 12, 11, 6–8] and the combinatorics of the strictly convex sets [4]. In the general case, the digitization set can be seen as the consequence of a group action on the object. A function, so-called *dual*, linking the group action and the produced digitization is used by the authors to study the digitization set up to a translation, for function graphs in [1] and for planar object in [10]. In the latter case, the dual has been proved to be piecewise constant in function of the translation. This paper focuses on the dual combinatorics.

Two upper bounds are given for the number of digitizations of a planar object whose boundary is a Jordan curve. The first one is expressed in terms of the number of grid cells crossed by the boundary and the second one in terms of the intersection number when plotting the boundary on the torus  $\mathbb{R}^2/\mathbb{Z}^2$ . The latter bound is proved to be quadratic in the convex case. Some examples are provided in order to compare the two upper bounds both in the convex and the non-convex cases. A conclusion and some perspectives end the paper.

## 2 Background

Let us consider a connected compact set  $S$  in  $\mathbb{R}^2$  whose boundary is a simple closed (Jordan) curve  $\Gamma$ . Thanks to the Jordan curve theorem, we may assume a continuous map  $f: \mathbb{R}^2 \rightarrow \mathbb{R}$  such that  $\Gamma$ , resp.  $S$ , is implicitly defined by

$$\Gamma = \{f(x) = 0 \mid x \in \mathbb{R}^2\} ,$$

resp.

$$S = \{x \in \mathbb{R}^2 \mid f(x) \leq 0\} .$$

We are interested in the variability of the Gauss digitization when the group of the translations acts on  $S$ , that is on the sets  $(u + S) \cap \mathbb{Z}^2$ ,  $u \in \mathbb{R}^2$ . In this paper, we focus on the combinatorial aspects of this variance. Of course, the variance has to be understood “up to integer translations”. This is the reason why we defined in a previous paper [10] the *dual by translation* of the digitizations of  $S$  as a set-valued function  $\varphi_S$  defined on the torus  $\mathbb{T} = \mathbb{R}^2 / \mathbb{Z}^2$  which maps each point  $t \in \mathbb{T}$  to the digitization, up to integer translations, of  $u + S$  where the vector  $u$  is any representative of  $t$  in  $\mathbb{R}^2$ . Let us pick a representative in each class of  $\mathbb{T}$  so as to form a connected set  $M$  called the *structuring element*. We note  $C$  the symmetric of  $M$  with respect to the origin:  $C = -M$ . The family of sets  $p + C$ ,  $p \in \mathbb{Z}^2$ , is a tiling of  $\mathbb{R}^2$ . For instance, we can take  $M = [0, 1]^2$ ,  $C = (-1, 0]^2$ . From now on, to simplify the notations, we confuse  $t \in T$  with its representative in  $M$  and the subsets of  $\mathbb{Z}^2$  with their orbits for the action of the integer translations so we can write  $\varphi_S(t) = (t + S) \cap \mathbb{Z}^2$ . For any point or set  $X$  in  $\mathbb{R}^2$ , we denote by  $\text{proj}(X)$  its projection on the quotient space  $\mathbb{T}$ .

We define the *grid boundary*  $\mathcal{B}$  as the set of grid points that lie in the (morphological) dilation of the boundary  $\Gamma$  of  $S$  by the structuring element  $M$ :

$$\mathcal{B} = (\Gamma \oplus M) \cap \mathbb{Z}^2 ,$$

where  $\oplus$  denotes the Minkowski sum. The set  $\mathcal{B}$  contains all the points of  $\mathbb{Z}^2$  whose membership to the digitization may change when the set  $S$  is shifted by a vector  $u \in M$ . Nevertheless, since  $\mathbb{R}^2$  is connected,  $M$  is not open so there may exist some points in  $\mathcal{B} \cap S$  not liable to change, namely those points  $p$  in  $\mathcal{B}$  for which  $p + C \subseteq S$  (see Fig. 1). That is why we have in fact to consider the *toggling boundary*  $\mathfrak{B}$  as the set of grid points whose membership effectively “toggle” for some translation by a vector  $u \in M$ :

$$\mathfrak{B} = \mathcal{B} \setminus \{p \in \mathbb{Z}^2 \mid p + C \subseteq S\} .$$

The set  $S \cap \mathbb{Z}^2 \setminus \mathfrak{B}$  of the grid points that are in any digitization of  $u + S$ ,  $u \in M$ , is called the *digitization core*.

Finally, for any  $p \in \mathbb{Z}^2$ , we set

$$\Gamma_p = ((-p) + \Gamma) \cap C ,$$

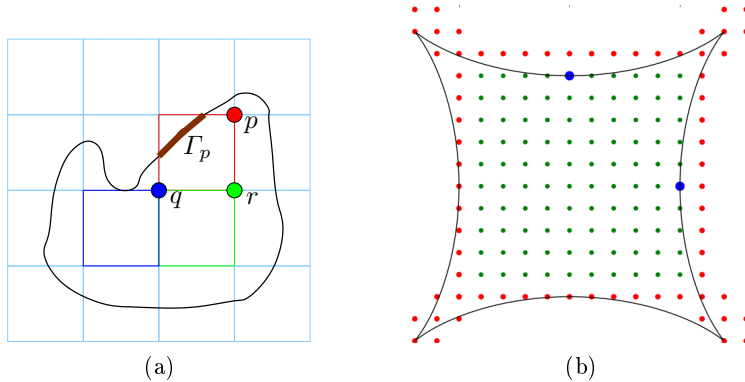


Fig. 1: (a) Black: a Jordan curve  $\Gamma$ ; light blue: tiles  $z + C$ ,  $z \in \mathbb{Z}^2$ ; red disc: a point  $p \in \mathfrak{B}$  (the red square  $p + C$  intersects  $\Gamma$  — in brown:  $\Gamma_p = (p + C) \cap \Gamma$  — and also intersects  $\mathbb{Z}^2 \setminus S$ ); blue disc: a point  $q \in \mathcal{B} \setminus \mathfrak{B}$  (the blue square  $q + C$  intersects the boundary  $\Gamma$  but is included in  $S$ ); green disc: a point  $r$  in the digitization core which is not in  $\mathcal{B}$  (the green square  $r + C$  is included in  $S$  and does not intersect  $\Gamma$ ). (b) Black: a Jordan curve, red: the toggling boundary, blue: two points that lie in the grid boundary but that cannot “toggle”, green: the digitization core.

so  $\Gamma_p = \emptyset$  if  $p \notin \mathcal{B}$ , and we denote by  $\mathbb{1}_p$  the indicator function of the set  $\Gamma_p$ , so  $\mathbb{1}_p$  is not constant iff  $p \in \mathfrak{B}$ . Then,

$$\varphi_S(t) = \{p \in \mathbb{Z}^2 \mid \mathbb{1}_p(-t) = 1\} .$$

We proved in [10] that the plot of the curve  $\Gamma$  on the torus  $\mathbb{T}$ ,  $\text{proj}(\Gamma)$ , delineates regions on which the dual function  $\varphi_S$  is constant.

*Property 1 ([10]).* Let  $S$  be a compact subset of  $\mathbb{R}^2$  whose boundary  $\Gamma$  is a Jordan curve. Then,

- the dual  $\varphi_S$  is constant on the connected components of  $\mathbb{T} \setminus \text{proj}(\Gamma)$ ;
- when  $M = [0, 1]^2$ , putting  $\mathbf{i} = (1, 0)$ ,  $\mathbf{j} = (0, 1)$ , it holds for any  $p \in \mathbb{Z}^2 \setminus \Gamma$  and any  $x, y \in [0, 1)$  that,
  - $\mathbb{1}_p(x\mathbf{i}) = \lim_{\lambda \rightarrow 1} \mathbb{1}_{p-\mathbf{j}}(x\mathbf{i} + \lambda\mathbf{j})$ ,
  - $\mathbb{1}_p(y\mathbf{j}) = \lim_{\lambda \rightarrow 1} \mathbb{1}_{p-\mathbf{i}}(y\mathbf{j} + \lambda\mathbf{i})$ .

In the rest of the article, we use the notation  $|E|$  for the cardinal of a set  $E$  ( $|E| \in \mathbb{N} \cup \{\infty\}$ ),  $\text{CC}(A)$  for the family of the connected components of a subset  $A$  of  $\mathbb{R}^2$  and  $I \sqcup J$  for the disjoint union of the sets  $I$  and  $J$ .

In the following section, upper bounds are given for the number of digitizations up to a translation. The first upper bound is naive. For the second one, the idea is to link the digitization number to the number of the intersection of the translated curves  $\Gamma_p$ . Indeed, the dual induces a partition of the torus whose cells frontiers are arcs of the curve  $\text{proj}(\Gamma)$ . Then, we bound from above the partition size by counting the number of curve intersections in  $\text{proj}(\Gamma)$ .

### 3 Contribution

#### 3.1 Bounding up by counting the crossed tiles

Since the grid boundary  $\mathcal{B}$  contains any integer point whose value may change when shifting the set  $S$ , we have a first, obvious, upper bound on the number of Gauss digitizations given by  $2^{|\mathcal{B}|}$ . As  $|\mathcal{B}|$  is also the number of tiles  $p + C$ ,  $p \in \mathbb{Z}^2$ , crossed by the frontier of  $S$ , we can state the following proposition.

**Proposition 1.** *The number of Gauss digitizations, up to integer translations, is upper bounded by  $2^{a(\Gamma)}$  where  $a(\Gamma)$  is the number of tiles crossed by  $\Gamma$ .*

Generally, the digitization enumeration provided by Prop. 1 includes false positives and multiple counts. For instance, the grid boundary of a circle with diameter 1.7 involves 4 to 8 pixels, depending on the grid position, which gives an upper bound according to Prop. 1 equals to 16, while there exists only 8 digitizations (see Fig 2). Nevertheless, from any set  $S$ , it is possible to build a new set that avoid false positives (but not multiple counts) by replacing the initial boundary by a family of Hilbert curves (see Appendix B).

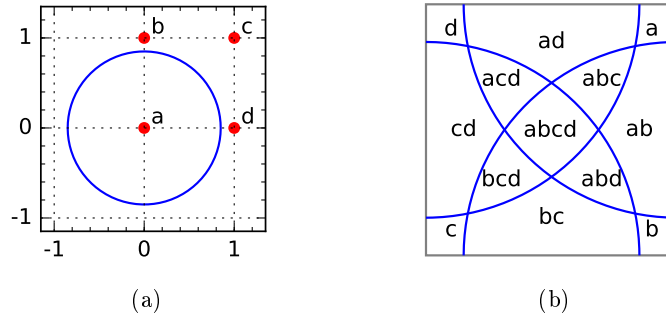


Fig. 2: (a) A circle with diameter 1.7 and the four pixels in the toggling boundary which is equal to the grid boundary: **a**, **b**, **c**, **d**. (b) The (flatten) dual of the closed disc bounded by the circle. Among the  $16 = 2^4$  potential digitizations provided by the 4-pixels **a**, **b**, **c**, **d**, three of them does not actually appear ( $\emptyset$ , **ac** and **bd**) and some others appear multiple times (the four singular digitizations, which are obviously congruent, and the vertical and horizontal pairs which each appear twice). Eventually, there are only 8 digitizations (and actually the (sewed) dual is divided into 8 regions).

#### 3.2 Bounding up by counting the intersections

In Sec. 3.2, we assume a parametrization of  $\Gamma$ . It induces an order on the points of the curve  $\Gamma$  (for  $\Gamma$  is simple) that is used in the proof of the following proposition.

The dual  $\varphi_S$  can be regarded as the projection, on the torus  $\mathbb{T}$  of a finite labeled partition of the tile  $C$  (whose cells need not be connected). **For the order by refinement on the partitions**, this partition is lower bounded by the infimum of the binary partitions  $\mathcal{P}_p$  associated to the indicator functions  $\mathbb{1}_p$ ,  $p \in \mathfrak{B}$  (actually,  $\text{proj}(\mathbb{1}_p)$ ,  $p \in \mathfrak{B}$ ). Indeed some cells of  $\bigwedge_{p \in \mathfrak{B}} \mathcal{P}_p$  need to be merged whenever the corresponding digitizations are equivalent up to an integer translation. Then, the next proposition proposes an upper bound on the size of the partition  $\bigwedge_{p \in \mathfrak{B}} \mathcal{P}_p$  — and thereby an upper bound to the number of digitizations — by inductively counting the intersections between the curves  $\Gamma_p$ ,  $p \in \mathfrak{B}$ . The idea is to count the partition cells created when adding a curve  $\Gamma_p$ . To do this we count the intersection of  $\Gamma_p$  with the already added curves. Nevertheless, these intersections can be with one or more curves and can be a singleton, **several points or an arc**. Then, let us define what means intersection in this context.

Definition 1 is illustrated in Fig. 3.

**Definition 1.** Let  $\mathfrak{B} = \{b_1, \dots, b_i, \dots, b_n\}$ ,  $n \geq 1$ . Let  $m \in [2, n]$ . Then,

$$\begin{aligned} \text{inter}_{\gamma, I} &= \text{CC} \left( \gamma \cap \left( \bigcap_{i \in I} \Gamma_{b_i} \setminus \bigcup_{j \in J} \Gamma_{b_j} \right) \right) \\ &\quad \text{where } \gamma \in \text{CC}(\Gamma_{b_m}) \text{ and } I \sqcup J = [1, m-1] \text{ ,} \\ \# \text{inter}_m &= \sum_{\substack{\gamma \in \text{CC}(\Gamma_{b_m}) \\ \emptyset \subset I \subseteq [1, m-1]}} w_I |\text{inter}_{\gamma, I}| \quad \text{where } w_I = \min(|I|, 2) \text{ .} \end{aligned}$$

*Note that the components of all the  $\text{inter}_{\gamma, I}$  are two by two disjoint and that  $\# \text{inter}_m$  may be infinite.*

The set  $\text{inter}_{\gamma, I}$  stands for the intersection of the curve  $\gamma$  with the curves whose indexes are in a subset  $I$ , excluding any other curve  $\Gamma_{b_i}$  whose index is not in  $I$ . The necessity to have all the index subsets comes from the fact that intersections between multiple curves can occur.

We are now able to state the proposition that relates the number of digitizations and the number of intersections  $\# \text{inter}_m$ .

**Proposition 2.** *The size of the partition  $\bigwedge_{p \in \mathfrak{B}} \mathcal{P}_p$  is upper bounded by*

$$2 + \sum_{m=2}^n \# \text{inter}_m + |\text{CC}(\Gamma_{b_m})| \text{ .}$$

*Proof.* We give here the sketch of the proof. The complete proof, which is tedious due to the large configuration space for the intersection set, can be found in the extended version of the paper [?].

The proposition is proved by induction. For  $m = 1$ , the result is obvious since, for any  $p \in \mathfrak{B}$ ,  $\mathcal{P}_p$  is a binary partition. Let  $m > 1$ . We assume that the number of cells of the partition  $\bigwedge_{i=1}^{m-1} \mathcal{P}_{b_i}$  is upper bounded by  $2 + \sum_{k=2}^{m-1} \# \text{inter}_k + |\text{CC}(\Gamma_{b_k})|$ . The cells of  $\bigwedge_{i=1}^{m-1} \mathcal{P}_{b_i}$  which are included in one

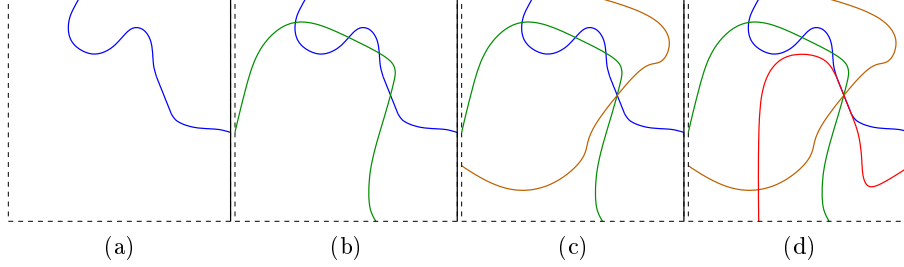


Fig. 3: (a) Blue: the curve  $\Gamma_{b_1}$ . (b) Green:  $\Gamma_{b_2}$  ( $\# \text{inter}_2 = |\text{inter}_{\Gamma_{b_2}, \{1\}}| = 4$ ). (c) Brown:  $\Gamma_{b_3}$  ( $\# \text{inter}_3 = |\text{inter}_{\Gamma_{b_3}, \{1\}}| + |\text{inter}_{\Gamma_{b_3}, \{2\}}| + 2|\text{inter}_{\Gamma_{b_3}, \{1,2\}}| = 0 + 0 + 2 = 2$ ). (d) Red:  $\Gamma_{b_4}$  ( $\# \text{inter}_4 = |\text{inter}_{\Gamma_{b_4}, \{1\}}| + 2|\text{inter}_{\Gamma_{b_4}, \{1,2,3\}}| = 2 + 2 = 4$ ).

of the two cells of  $\mathcal{P}_{b_m}$ , namely the sets  $(\mathbb{1}_{b_m} = 0)$  and  $(\mathbb{1}_{b_m} = 1)$ , stay unchanged in the partition  $\bigwedge_{i=1}^m \mathcal{P}_{b_i}$ . Conversely, the cells of  $\bigwedge_{i=1}^{m-1} \mathcal{P}_{b_i}$  that are intersected by both  $(\mathbb{1}_{b_m} = 0)$  and  $(\mathbb{1}_{b_m} = 1)$ , which obviously are cells intersected by  $\Gamma_{b_m}$ , are each divided in two new cells. Thereby, the number  $N_m$  of new cells is upper bounded by  $N'_m$ , the number of cells in  $\bigwedge_{i=1}^{m-1} \mathcal{P}_{b_i}$  intersected by  $\Gamma_{b_m}$ . Besides, the partition  $\bigwedge_{i=1}^{m-1} \mathcal{P}_{b_i}$  of the tile  $C$  induces a partition  $\mathcal{Q}$  of  $\Gamma_{b_m}$ , as a subset of  $C$ , with  $N'_m$  cells (these cells of  $\Gamma_{b_m}$  need not be connected). Then, the idea of the proof is to map each cell of  $\mathcal{Q}$  to its supremum, for the order induced by the parametrization of  $\Gamma$  — namely to the intersection, as defined in Def. 1 — its supremum belongs to, or to the empty set when no such intersection exists. Then, though this mapping is not one-to-one, a careful examination of the different cases permits us to conclude that  $N'_m$  is upper bounded by  $\# \text{inter}_m + |\text{CC}(\Gamma_{b_m})|$ .  $\square$

When projecting the partition  $\bigwedge_{i=1}^n \mathcal{P}_{b_i}$  on the torus  $\mathbb{T}$ , the cells that touch the boundary of  $C$  are identified two by two, which decreases the number of cells in the partition. Unfortunately, it is difficult to count these cells in the general case. Nevertheless, in the most simple case (each curve crosses the boundary of  $C$  in two points, all points and cells are distinct), the reader can easily check that the identification leads to remove the term  $\sum_{i=1}^n |\text{CC}(\Gamma_{b_i})|$  in Prop. 2.

### 3.3 The convex case

It seems plain that the structure of the dual should be simpler when the set  $S$  is convex compare to a set with a winding boundary. We could even hope that the digitizations coincide with the cells of the partition  $\bigwedge_{p \in \mathfrak{B}} \Gamma_p$ , after sewing the boundary of the tile  $C$ , and with the connected regions of the torus  $\mathbb{T}$  delineated by the curve  $\Gamma$ . Figure 4 annihilates this hope by exhibiting a convex object and one of its digitizations whose inverse image by the dual is not connected.

Nevertheless, in this section we show that the structure of the sets  $\text{inter}_{\gamma, I}$  is simple when the set  $S$  is strictly convex and permits to obtain a quadratic

bound for the complexity of the dual in terms of the grid boundary size (to be compared to the exponential bound of Prop. 1).

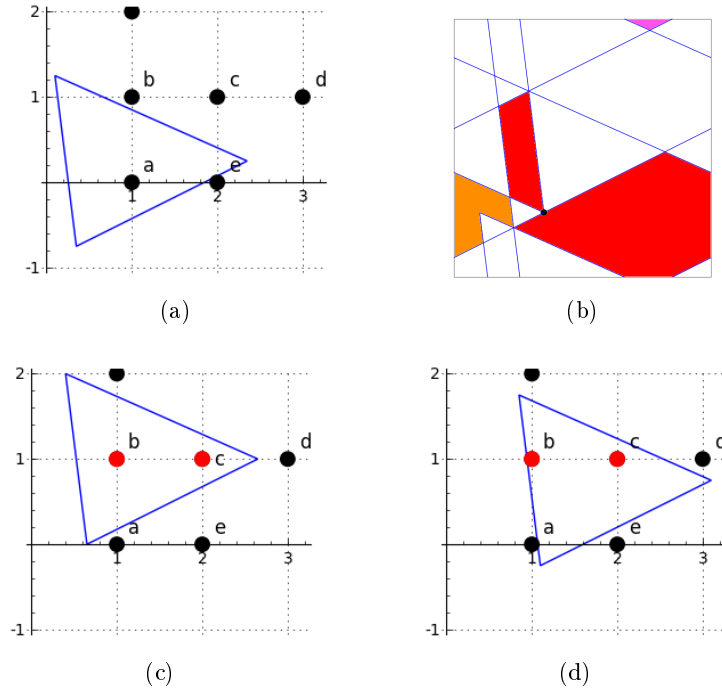


Fig. 4: (a) A triangle and its toggling boundary. (b) The dual of the (filled) triangle. The four colored regions correspond to the same digitization: the horizontal pair. Moreover, the two red regions belong to the same cell of the partition  $\bigwedge_{p \in \mathfrak{B}} \Gamma_p$ , the cell associated to the digitization **bc** (these two regions are separated by the unique point of the cell **abcd**). The orange region is the cell **cd** and the pink region is the cell **ae**. (c), (d) The configurations of the shifted triangle in the two red regions. Since they are disconnected, it is not possible to continuously move to one configuration to the other without hurting a black toggling point. Note that the triangle could be slightly inflated so as to provide a strictly convex counterexample.

**Proposition 3.** *We assume a convex quadrilateral structuring element  $M$ . Let  $n = |\mathfrak{B}|$ . The number of Gauss digitizations of a strictly convex planar object up to a translation is upper bounded by*

$$4n^2 + 4n - 6 .$$

*Proof.* The proof invokes Lemma 3 which is stated and proved in the appendix. Let  $N$  be the digitization number. Thanks to Prop. 2, we have

$$N \leq 2 + \sum_{m=2}^n \# \text{inter}_m + |\text{CC}(\Gamma_{b_m})| . \quad (1)$$

From Lemma 3, we derive that, for any  $m \in [1, n]$ ,  $\gamma \in \text{CC}(\Gamma_{b_m})$ ,  $i \in [1, m-1]$ ,

$$\sum_{I \ni i} |\text{inter}_{\gamma, I}| \leq 2 |\text{CC}(\Gamma_{b_i})| .$$

Then,

$$\# \text{inter}_m \leq 2 \sum_{k=1}^{m-1} |\text{CC}(\Gamma_{b_k})| .$$

Eventually, Eq (1) turns into

$$\begin{aligned} N &\leq 2 + 2 \sum_{m=2}^n \sum_{k=1}^{m-1} |\text{CC}(\Gamma_{b_k})| + \sum_{m=2}^n |\text{CC}(\Gamma_{b_m})| \\ &\leq 2 \left( 1 + \sum_{k=1}^n (n-k+1) |\text{CC}(\Gamma_{b_k})| - |\text{CC}(\Gamma_{b_1})| \right) . \end{aligned}$$

It can easily be seen (from Lemma 2 for instance) that, for any convex curve  $\Gamma$  and any convex polygon  $P$  with  $c$  edges, the number of connected arcs in the intersection  $\Gamma \cap P$  is upper bounded by  $c$ . Then, if we assume that the tile  $C$  is a convex quadrilateral, we straightforwardly obtain the desired bound.

The term  $4n^2$  in Prop. 3 comes down to  $n^2$  for sufficiently high resolution because each curve  $\Gamma_{b_i}$  will have just one connected component instead of possibly 4 in the general case. Then, the result of Prop. 3 is close to the one in [6] which states that the number of digitizations of a disc of radius  $r$  is asymptotically  $4\pi r^2 + O(r)$ . As the ratio between the radius  $r$  and the size of the grid boundary is  $\pi/4$  for the disc, our upper bound in function of the radius  $r$  of the disc is asymptotically equivalent to  $16/\pi^2 r^2$ .

## 4 Conclusion

We present in this paper two upper bounds on the number of digitizations obtained from all the translated of a continuous object. The first one is exponential in the number of toggling object-boundary pixel and a generic example reaching this bound is given. The second one is based on the passage from the dual connected-component count to the curve-intersection count. If the curve is parametrized, then the count comes down to count some equation solutions. The second upper bound is then used in the convex case, where digitization number is shown to be quadratic. An example of a convex object is given where the



set of translation parameter classes corresponding to a given digitization is not connected.

The perspectives of this work are first to explicit the second upper bound under assumptions less restrictive than the convexity, *e.g.* bounded curvature; then to study the combinatorics of the digitization under the rigid transformations and to propose an algorithm for the digitization generation.

## A Convex sets

The proof of Lemma 3 relies on the two following lemmas about convex sets that seem obvious at the first glance. Nevertheless, since we did not find any result related to these lemmas in the literature, we provide our own justifications of the two statements.

**Lemma 1 (Chords of convex sets).** *Let  $[a, b]$  be a chord of the boundary  $\Gamma$  of a closed convex set  $S$ . If  $[a, b] \not\subseteq \Gamma$ , then  $(ab) \cap S = [a, b]$  and  $(ab) \cap \Gamma = \{a, b\}$ .*

*Proof.* Since  $[a, b] \not\subseteq \Gamma$ , the line  $(ab)$  does not support  $S$  at any point (notion de supporting half-plane : Berger ou Valentine). So, there exists two supporting lines of  $S$  at  $a$  and  $b$  that cross the line  $(ab)$ . Then,  $(ab) \cap S$  is included in  $[a, b]$ . Let  $c \in [a, b] \cap \Gamma$ ,  $c \neq a$ . Applying the first part of the proof to the chord  $[a, c]$ , we derive that  $(ac) \cap S \subseteq [a, c]$  and, since  $b \in (ac) \cap S$ , we conclude that  $b = c$ .

**Lemma 2 (Cuts of convex sets).** *Let  $a, b$  be two points of the boundary  $\Gamma$  of a closed convex set  $S$ . If  $[a, b] \not\subseteq \Gamma$ , then the open curve segments of  $\Gamma$ ,  $\sigma_1, \sigma_2$ , whose extremities are  $a$  and  $b$  are included in distinct open half-planes bounded by the line  $(ab)$ .*

*Proof.* Let  $H^-$  and  $H^+$  be the two open half-planes bounded by  $(ab)$ . Since  $[a, b] \not\subseteq \Gamma$ , from Lemma 1,  $(ab) \cap \Gamma = \{a, b\}$ . Thereby, by connectivity, either  $\sigma_1 \subset H^-$  or  $\sigma_1 \subset H^+$  and  $\sigma_2 \subset H^-$  or  $\sigma_2 \subset H^+$ . Suppose for instance that  $\sigma_1 \subset H^-$  and  $\sigma_2 \subset H^-$ . Then,  $S$ , which is the connected subset of  $\mathbb{R}^2$  bounded by  $\sigma_1 \cup \sigma_2 \cup \{a, b\}$  (citeJordan) is included in  $H^- \cup (ab)$  and, since  $(ab) \cap S = [a, b]$  from Lemma 1,  $[a, b] \subset \Gamma$ : contradiction.

**Lemma 3 (Intersection of two segments of a convex curve).** *Let  $\Gamma$  be a Jordan curve whose interior is convex. Let  $\Gamma_1$  and  $\Gamma_2$  two disjoint (closed) segments/arcs of the curve  $\Gamma$  and  $\tau$  a vector of  $\mathbb{R}^2$ . Then the intersection of  $\Gamma_1$  and  $\tau + \Gamma_2$  is composed of none, one, two points or a line segment.*

*Proof.* Let  $p, q$  be two distinct points in  $\Gamma_1 \cap (\tau + \Gamma_2)$  if such a pair exists. We denote by  $\Sigma_1$  the open segment of  $\Gamma_1$  between  $p$  and  $q$ . Alike,  $\Sigma_2$  is the open segment of  $\Gamma_2$  between  $-\tau + p$  and  $-\tau + q$ . We set  $\overline{\Sigma}_1 = \Sigma_1 \cup \{p, q\}$  and  $\overline{\Sigma}_2 = \Sigma_2 \cup \{p, q\}$ . Firstly, we prove that  $\Sigma_1 \cup (\tau + \Sigma_2)$  is a straight line segment whenever it contains more than two points.

First case:  $\Sigma_1 \cup (\tau + \Sigma_2) \subseteq (pq)$ . Then, since  $\overline{\Sigma}_1$  and  $\overline{\Sigma}_2$  are connected and  $\Gamma$  is simple,  $\overline{\Sigma}_1 = \overline{\Sigma}_2 = [p, q]$ .

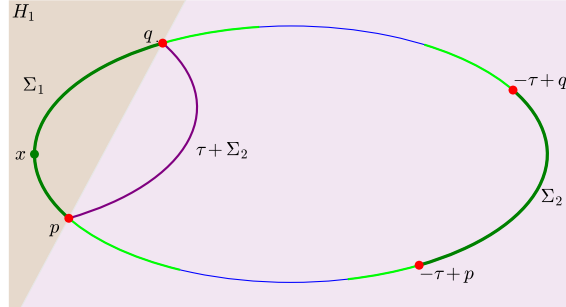


Fig. 5

Second case:  $\exists x \in (\Sigma_1 \cup (\tau + \Sigma_2)) \setminus (pq)$ . For instance, we assume  $x \in \Sigma_1 \setminus (pq)$ . By Lemma 2,  $\Sigma_1 = \Gamma \cap H_1$ , where  $H_1$  is the open half-plane bounded by the line  $(pq)$  and containing  $x$ , and  $-\tau + p$  is in  $\mathbb{R}^2 \setminus H_1$ . Then, it can easily be seen that  $-\tau + H_1$  is the open half-plane bounded by the line joining  $-\tau + p$  and  $-\tau + q$  and including  $p$ . Thanks to Lemma 2, we derive that  $\Sigma_2$  does not intersect  $-\tau + H_1$ . Thus,  $\tau + \Sigma_2$  does not intersect  $H_1$ . In particular,  $(\tau + \Sigma_2) \cap \Sigma_1 = \emptyset$ . This achieves the first part of the proof.

Now, let  $r$  be a point in  $\Gamma_1 \cap (\tau + \Gamma_2)$  which is not in  $\Sigma_1$  (if such a point exists). For instance,  $p$  belongs to the segment of  $\Gamma_1$  between  $q$  and  $r$ . Then, the first part of the proof, applied to the points  $q$  and  $r$ , implies that  $\Gamma_1 \cap (\tau + \Gamma_2)$  includes the segments  $[q, r]$ . We straightforwardly concludes that either the intersection of  $\Gamma_1$  and  $\tau + \Gamma_2$  is composed of at most two points or it is a line segment.

## B Examples and counterexamples

### B.1 Building examples without proper congruent digitizations in the image of the dual

Let  $u$  and  $v$  be two vectors in  $[0, 1)^2$  such that the sets  $u+S$  and  $v+S$  have distinct but congruent digitizations. Then, there exists an integer vector  $w$ ,  $w \neq 0$ , such that  $(u+S) \cap \mathbb{Z}^2 = w + ((v+S) \cap \mathbb{Z}^2) = (w+v+S) \cap \mathbb{Z}^2$ . Let  $p$  be a point in the digitization core. Then,  $p \in (u+S) \cap \mathbb{Z}^2$  and  $p \in (v+S) \cap \mathbb{Z}^2$ . Therefore,  $w+p \in (u+S) \cap \mathbb{Z}^2$  and  $-w+p \in (v+S) \cap \mathbb{Z}^2$ , which can be rewritten as  $p \in ((-w+u)+S) \cap \mathbb{Z}^2$  and  $p \in ((w+v)+S) \cap \mathbb{Z}^2$ . Then, at least one of the vector  $w+v$  or  $-w+u$  has one of its coordinates which is negative. We derive that if there exists a point in the digitization core which is maximal in  $S$  for both coordinates then there is no proper congruent digitizations in the dual. An example with such a maximal point in the digitization core is provided in Fig. 6.

## B.2 Building toric partitions in one-to-one correspondence with the power set of the toggling boundary

In this section, we exhibit a way to modify the boundary of the set  $S$  in order to ensure that any set in the power set of the toggling boundary is represented in the dual. To do so, we move along  $\mathcal{B}$ , ordered in the same way as in Def. 1. Then, a new boundary is built thanks to the approximations of the Hilbert filling curve: the segment of  $L$  intersecting the  $n$ -th cell of  $\mathcal{B}$  is replaced by a  $n$ -th approximation of the Hilbert filling curve  $H_n$  (extended at its extremities to ensure the continuity of the boundary): see Fig. 6. We consider the family

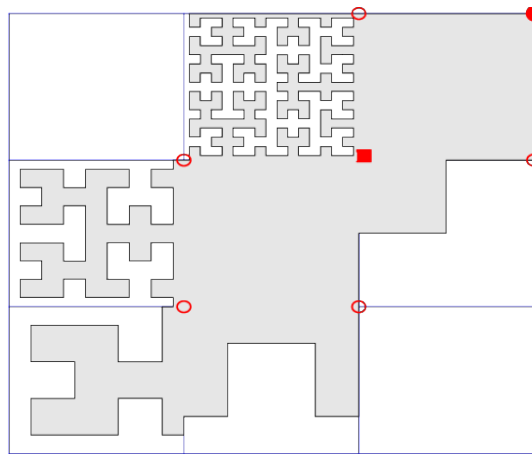


Fig. 6: The boundary of the set  $S$  depicted in the figure is obtained by connecting  $n$ -th approximations of the Hilbert filling curve,  $0 \leq n \leq 4$ . Red square and red disc: digitization core. Red circle: the five points in the toggling boundary. The red disc is both in the grid boundary and in the digitization core. Furthermore, it is maximal for both coordinates which ensures that there is no proper congruent digitizations in the image of the dual. The set  $S$  has  $2^5$  digitizations.

of binary partitions  $\mathcal{P}_n$  of the unit square that comes with the curves  $H_n$ . We claim that each curve  $H_{n+1}$  crosses each cell of the partition  $\bigwedge_{i=1}^n \mathcal{P}_i$  so that the size of the final torus partition is  $2^N$  where  $N$  is the cardinal of the toggling boundary. To justify our claim, we divide the unit square in a family of  $2^n \times 2^n$  small squares  $(K_{i,j}^n)_{1 \leq i,j \leq 2^n}$  ( $n \geq 0$ ) whose sizes are  $\frac{1}{2^n} \times \frac{1}{2^n}$ . It can be seen that on the one hand, the Hilbert curve  $H_n$  passes through the center of each of the squares  $K_{i,j}^n$  and, on the other hand, does not intersect any of the interior of the squares  $K_{i,j}^{n+1}$  ( $H_0$  is just the center of the unit square). Thereby, the partition  $\bigwedge_{i=1}^n \mathcal{P}_i$  is coarser than the partition  $\{K_{i,j}^{n+1} \mid 1 \leq i,j < n\}$  (the boundaries of the squares  $K_{i,j}^k$  are assigned to the cells so as to coincide with  $\mathcal{P}_n$ ). Since  $H_{n+1}$

passes through the center of each of the squares  $K_{i,j}^{n+1}$ , it passes in each cell of  $\bigwedge_{i=1}^n \mathcal{P}_i$  which gives the claim.

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